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Dipolar Fields at Great Distances with Hybrid Parallelism

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Abstract:

Because they grow more quickly with system size compared to the other interactions, calculations involving dipolar fields take up the bulk of the time in LLG models. For a parallel LLG technique, effectively parallelizing the dipolar field is crucial. The approach described in this piece to actively create a parallel plan and weigh its merits against more traditional forms of parallelization.

Introduction

The dipolar component to the effective field in the Landau-Lifshitz-Gilbert (LLG) equation is determined using the Fast Fourier Transform (FFT). This is what has drastically reduced the computational time required for micromagnetic models from O to O(N log(N)).(N2). Calculating The dipolar fields can be implemented in a parallel setting using standard parallel- alleviation techniques, such as coarse or fine grain parallelism. While these methods can work in parallel, they are rarely ideal due to the complexity of high-performance computing. It's important to note that the optimal solution for one problem size may not be applicable to another because the dipolar interaction and all other interactions scale differently in terms of time.Details of the dipolar calculation that make writing parallel code difficult will be presented below. We will also show how to instantly generate a scalable parallel strategy that takes into account the available resources, the number of concurrent users, the general scale of the system, and the scalability of the techniques being used.

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The Dipolar Interaction

The dipolar interaction is a long-range, low-intensity field that exists between all magnetic moments[1] and can help maintain order at room temperature[2]. The dipolar field's importance in the correct handling of thin films, despite its much lower size compared to the exchange interaction, magnetic thin film devices.

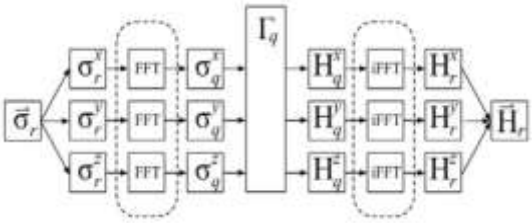

Figure 1: Schematic of Dipolar Field Calculation

The dipolar field at location I of a system with distinct lattice sites is given by

$$\vec{H}_i^{dd} = g\sum_j{}' \frac{\vec{\sigma}_j}{r_{ij}^3} - \frac{3\vec{r}_{ij}(\vec{\sigma}_j \cdot \vec{r}_{ij})}{r_{ij}^5}, \qquad (1)$$

To disprove the null hypothesis that I = j, we use the symbol j for the magnetic moment at position j, r in for the distance between places I and j, and a priming operation on the sum. If we characterize our global lattice as a single lattice that is endlessly repeated, we can disregard the irregular boundary condition repeats.

gravitational frameworks

$$(\vec{H}_i^{dd})^\alpha = g\sum_j \Gamma_{ij}^{\alpha\beta}\sigma_j^\beta, \qquad (2)$$

with interaction matrix $\Gamma^{\alpha\beta}$ as

$$\Gamma^{\alpha\beta} = \sum_{\vec{R}}{}' \frac{\delta^{\alpha\beta}}{|\vec{R}_{ij}|^3} - \frac{3R^\alpha R^\beta}{|\vec{R}_{ij}|^5}, \qquad (3)$$

which is a summation over replicated lattices $\vec{R}$. The Fourier transform of Equation (1) can now be written as

$$H^\alpha(\vec{q}) = \Gamma^{\alpha\beta}(\vec{q})\sigma^\beta(\vec{q}) \qquad (4)$$

which maps to a frequency-domain convolution of order N. The computations required to extract the dipolar fields in real space from a single layer of magnetic moments are summarized in Figure. (1). Indicated by these figures are Quick Fourier The magnetic moments must be transformed into frequency space and decomposed along Cartesian axes. Convolution with the previously calculated Fourier-transformed interaction matrices yields the dipolar fields in the centre. The produced areas must be transformed back to "regular" space after generation. This schematic illustrates how we perform several FFTs during each turn rather than just one. We need to parallelize the calculation of the dipolar fields so that it can be done quickly and accurately. We contend that the finest approach to coping with either of the two highlighted groups can be applied to evaluating the other, despite their differences.

Group Parallelization

Our goal is to create a method of parallelization that will facilitate the quick accomplishment of numbeuros tasks. In this case, Fast Fourier Transforms are employed. Several approaches have been proposed, and one of them makes use of coarse particulate parallelism, where the exe-Multiple machines work together to complete each task[3]. On the other hand, a thorough work scheduling and job parallelization to boost efficiency. Just thetraditional, time-tested processes. Coarse grain parallelism necessitates more CPUs than number of employment opportunities, which is a floor set too low. When sanding, high grain is ideal. It is also unfeasible because the counterpart expense is zero. Both of these methods are based on suppositions. The dynamic working environment and the extensiveness of the associated computations. constitute major flaws. Examples of such machines include the SGI Al- tera and other contemporary Symmetric Multi-Processor (SMP) systems. tix 350, have an inherently dispersed IPC delay structure. Two It is faster for data to travel between machines on the same node than between nodes. Comma- The time it takes for a message to move between nodes can also be affected by the method of transmission used by the nodes.employ a hypercube configuration, while others use a single planar high-velocity back plate. A It's possible that a hard-coded broad or fine-grained parallel response another. While it is feasible to write code for a single CPU, doing so for a collection of computers is not doable. on a single machine almost never does. Most SMP machines share hardware and software, dependence on all the time in the scholarly world. This introduces doubt into the optimal framework. It is in-processing cycle time. We have developed, tested, and released a solution to this problem. uses all of a computer's processors to calculate how well a battery of exams

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fared. This strategy has two parts. First, we select a set of idle CPUs and topologically closest, or very close. The set of tasks can be executed with the fewest restrictions using this choice. in a brisk, conversational fashion. The second stage involves formulating a plan that takes into think about how much work is currently being done on the system, what resources are available, how efficient the computations currently are, etc. We've decided to approach the problem in this way.1. We estimate how long it would take from 1 to N machines to finish the job before we even start the practice. Build a single task, an FFT, and label the final duration t.(n).2. We can divide our N processors into P pipelines, where in is the number of processors in each pipeline. [I = 1...P] The number of machines responsible for assessing the responsibilities of the itch line.3. For every in, we calculate all of the potential T task distributions. These are the P-traps. The number of tasks (denoted by It) that can be assigned to a given channel is then at your discretion. Pipelines Si = Nimit and the resulting system S = Si.

Conduit i's length is then specified as

$$t(S_i) = t(n_i)T_i, \tag{5}$$

$$t(\mathbf{S}) = \max(t(S_i)), \tag{6}$$

which indicates the overall amount of time required for all the pipelines to finish all the jobs given to them using the available CPUs. Pipes are done in simultaneously, with duties in each pipe being completed in order.

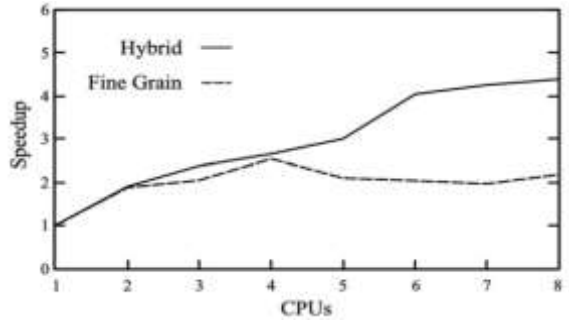

Figure 2: Comparison of Speedup for a 128x128 LLG simulation

The collection S I that minimizes t(S) subject to the restriction identifies the optimum number of pipelines, job allocation, and CPU count.

$$\sum_{i=1}^{P} n_i \leq N$$

$$\sum_{i=1}^{P} T_i = T. \tag{7}$$

Naturally, the collection of feasible solutions contains both fine and coarse granular plans. The collection corresponds to the fine-grained system.

$$\mathbf{S} = \{[N,T]\}, \tag{8}$$

while the coarse grained scheme coresponds to the solution

$$\mathbf{S} = \{[1,1],[1,1],\ldots,[1,1]\}, \quad |\mathbf{S}| = N. \tag{9}$$

At the outset of each simulation, we flexibly spread the T tasks across N CPUs to speed up the concurrent evaluation of the entire set of tasks. Because it utilizes both wide and fine grain parallelism, we call this approach hybrid parallelization. accomplish a top-notch result in the group assessment of concurrent activities.

Speedup:

In the context of a complete LLG simulation performed on the SGI Alix 350, Figure 2 provides an example contrast of the speedup achieved via the fine grain and mixed parallelism. This demonstrates that the combination approach identifies and chooses instances with fine grain parallelism. system and dual-processor system. Then, it picks the most effective strategies for situations with three or more processors. Both bigger and lesser grid sizes are represented by these shapes. Since waste eventually becomes dominant and lowers speedup, fine grain parallelism is only useful for a finite number of computers. As more processors are introduced, the hybrid parallelism changes from fine grain to coarse grain. It is possible that the best answer doesn't always use all of the CPUs. This is analogous to a conduit having zero jobs but an infinite amount of cores. If adding a CPU to the group causes more effort than it's worth, you'll run into this problem.

Conclusion

We have implemented a novel parallelism technique for FFT computation in groups. The dipolar interaction for an LLG simulation was then efficiently calculated using these FFTs. All possible parallel solutions between pure coarse and fine granular are taken into account by the technique. It evaluates the available parallelism techniques and picks the one that best suits the

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present processing environment and algorithmic efficiency. When compared to the fine granularity approach, this one can provide substantial speedup when expanding to more CPUs. The same holds true when contrasting with coarse grained parallelism, which is taken into account as a potential answer when an optimum plan is built. Hybrid parallelization strategies ensure performance outcomes that are on par with or better than those of conventional parallelization strategies because the collection of solutions evaluated by the strategy includes both coarse and fine granularity.

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